

Blockstream

Secure protocols on BIP-taproot

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GPG: 36C7 1A37 C9D9 88BD E825 08D9 B1A7 0E4F 8DCD 0366

Disclaimer

It's not at all certain that a BIP-taproot softfork activates in its current form or at all. This depends on community consensus.

BIP-taproot address generation (witness version 1)

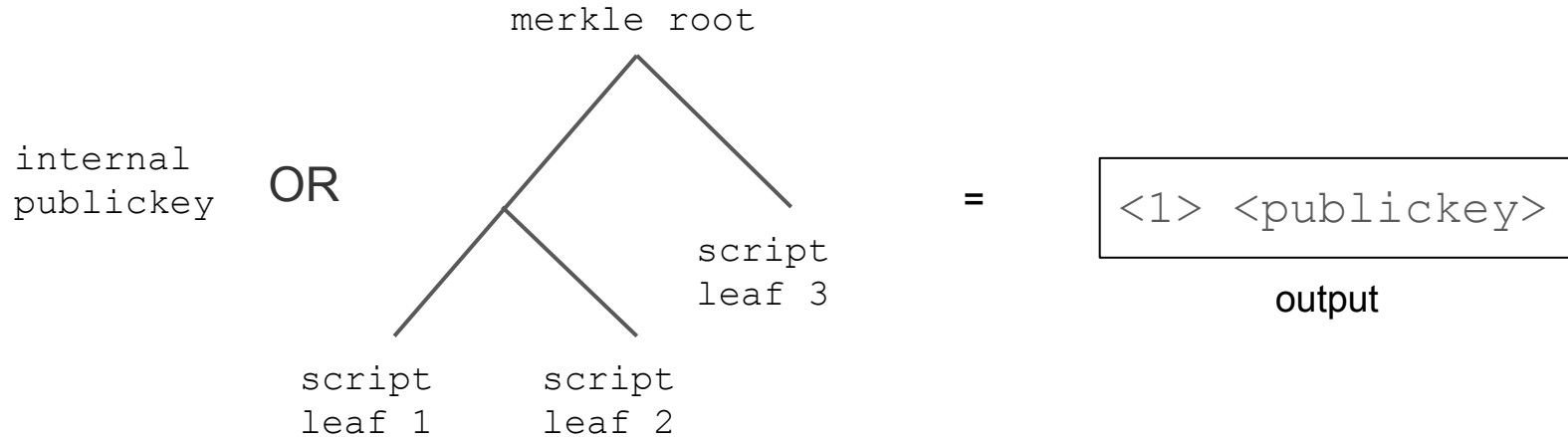
Policy: single key

```
<1> <publickey>
```

output

BIP-taproot address generation (witness version 1)

Policy: single key OR script1 OR script2 OR script3



BIP-taproot spending

output

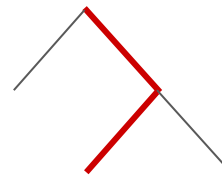
```
<1> <publickey>
```

Key spend

(BIP-schnorr) signature

Script spend (Script 2)

Script 2 inputs



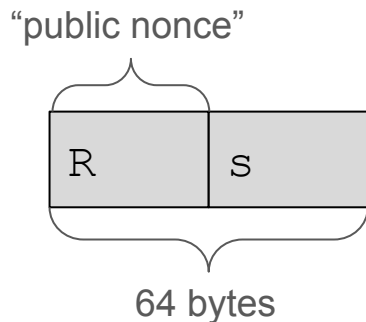
Script 2

bitcoin-core/secp256k1

- “Difficult to use insecurely”
 - Well reviewed and tested
 - Fast and portable
 - Free of timing sidechannels
- [rust-bitcoin/rust-secp256k1](#) type-safe rust bindings (`no_std`)
- Will provide cryptographic primitives for bip-taproot
 - minimum required: schnorrsig module
- [elementsproject/libsecp-zkp](#)
 - fork of secp256k1 with rangeproofs, surjectionproofs, schnorrsig, musig, ...
 - just released: [rust-secp256k1-zkp](#) beta (schnorrsig, optional `no_std`)
- HOWTO
 - read the docs before using it (`include/secp256k1_*.h`)

schnorr sig module

[secp256k1 PR #558](#)
[secp256k1-zkp module](#)
[rust-secp256k1-zkp module](#)

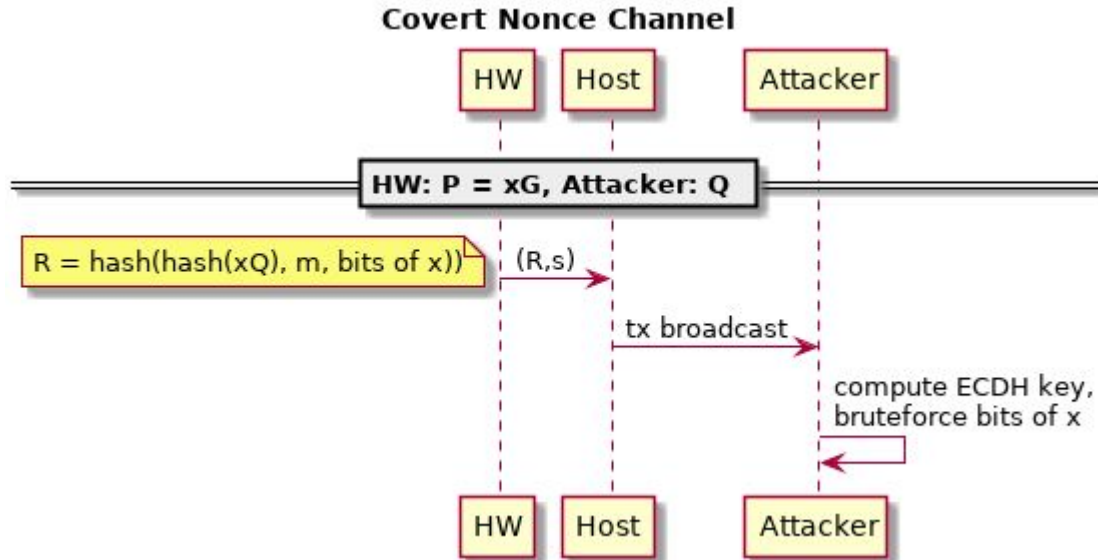


nonce = Number used ONCE

- Deterministic nonce derivation as per BIP-schnorr
 - Picking a specific nonce is unnecessary
- Batch verification
 - 400 sigs can be verified in half the time
 - Don't know which exact sig was invalid
 - May not reduce worst case cost

Covert nonce channel

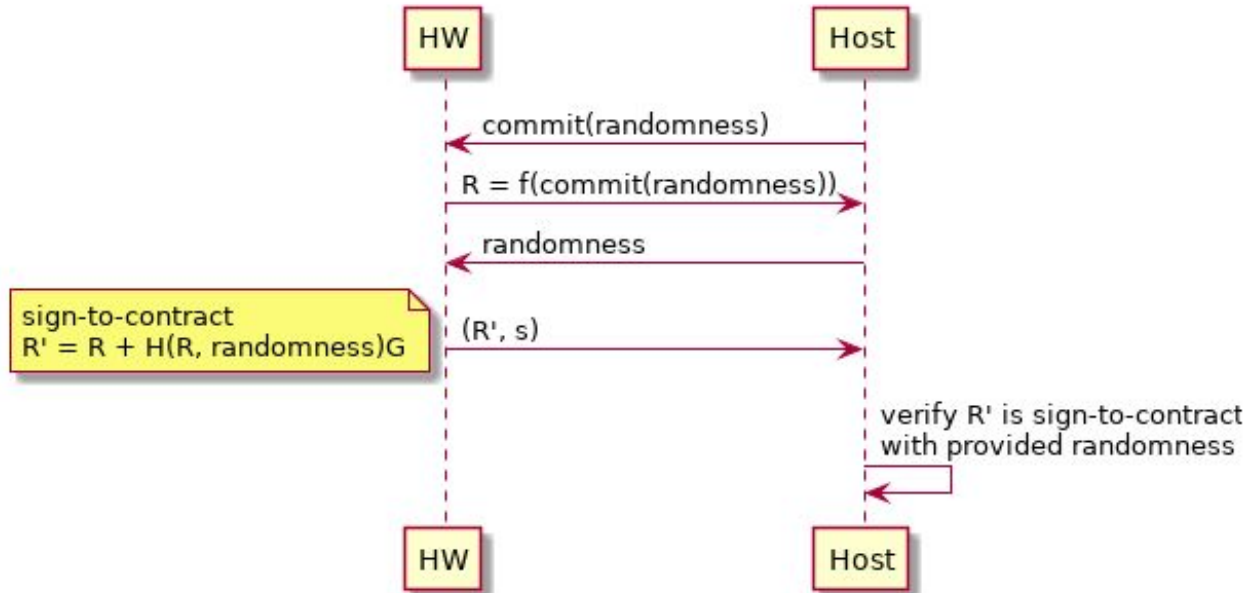
- Problem: malicious HWW can exfiltrate secret key through nonce



Covert nonce channel protection

Solution: enforce putting host-supplied randomness in nonce with *sign-to-contract*

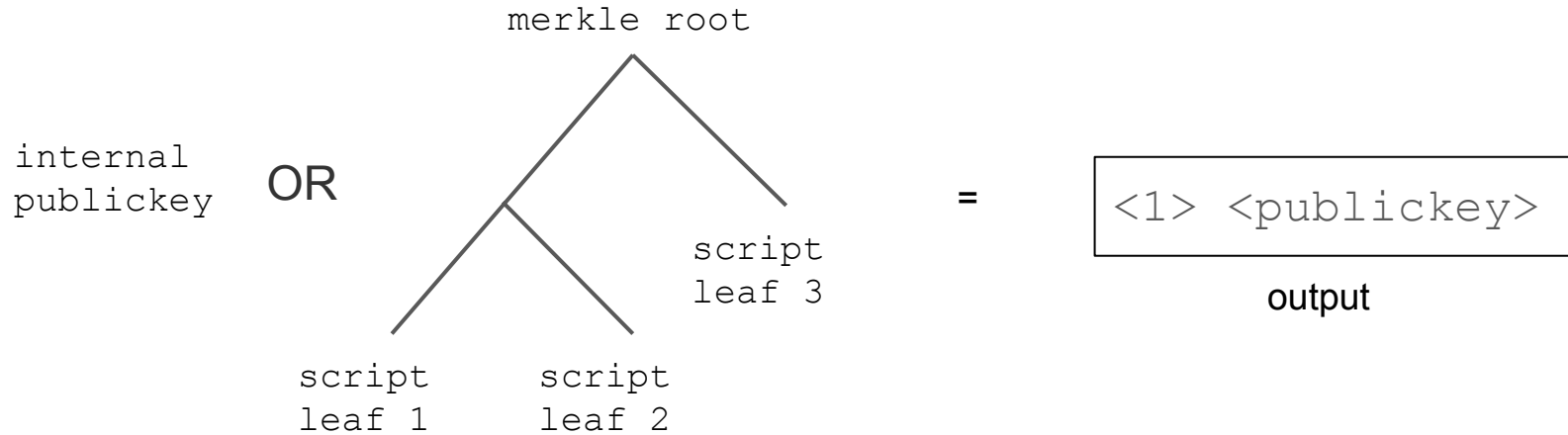
Covert Nonce Channel Protection



Alternative:
MuSig key
aggregation but
that's currently
difficult for
hardware wallets

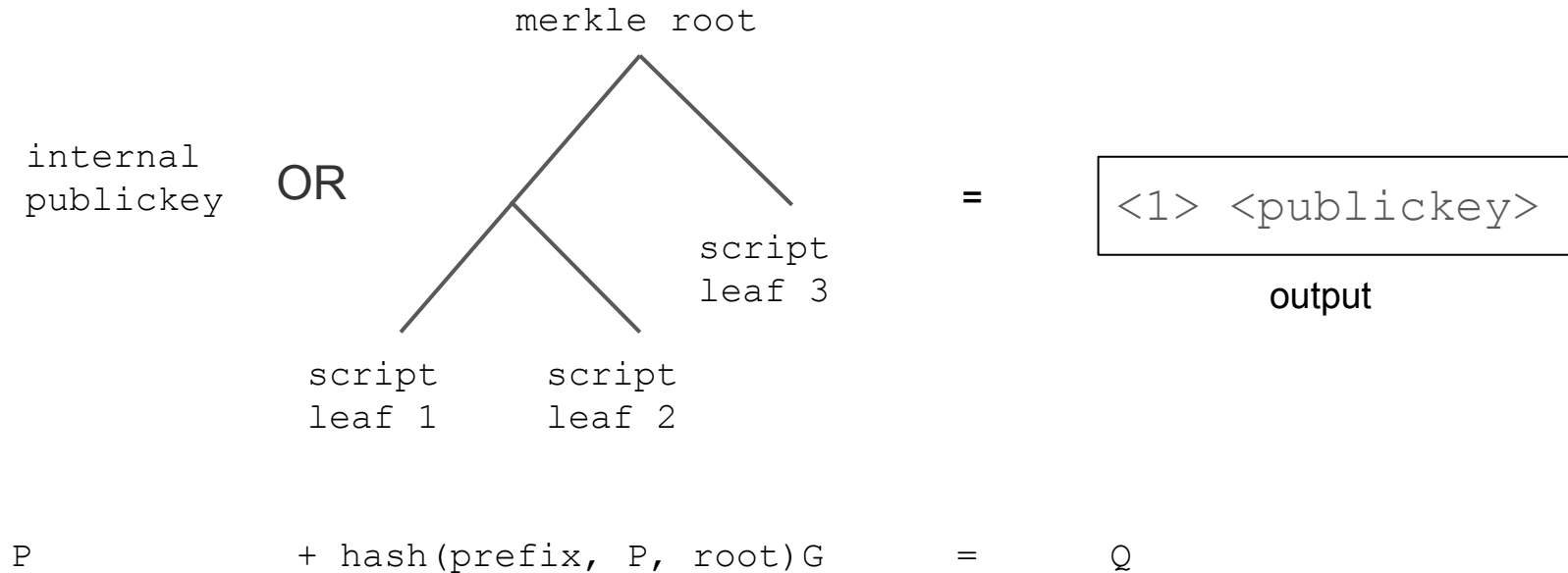
Tweak Add

- Create taproot commitment if there's a script path



Tweak Add

- Create taproot commitment if there's a script path



Tweak Add

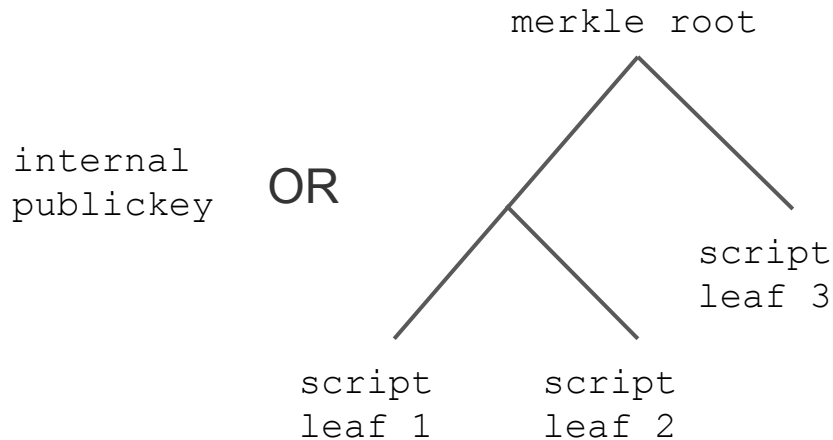
- Create taproot commitment if there's a script path

```
int secp256k1_ec_pubkey_tweak_add(  
    const secp256k1_context* ctx,  
    secp256k1_pubkey *pubkey,  
    const unsigned char *tweak)
```

$P + \text{hash}(\text{prefix}, P, \text{root})G = Q$

Tweak Add Fungibility

- Try avoiding the script path
 - in multi-party contracts use “happy” case
- Don't reuse keys
 - internal keys and leaf keys
- Using script path basically leaks wallet
 - Depth of tree, script, ...
- Ensure sufficient leaf entropy



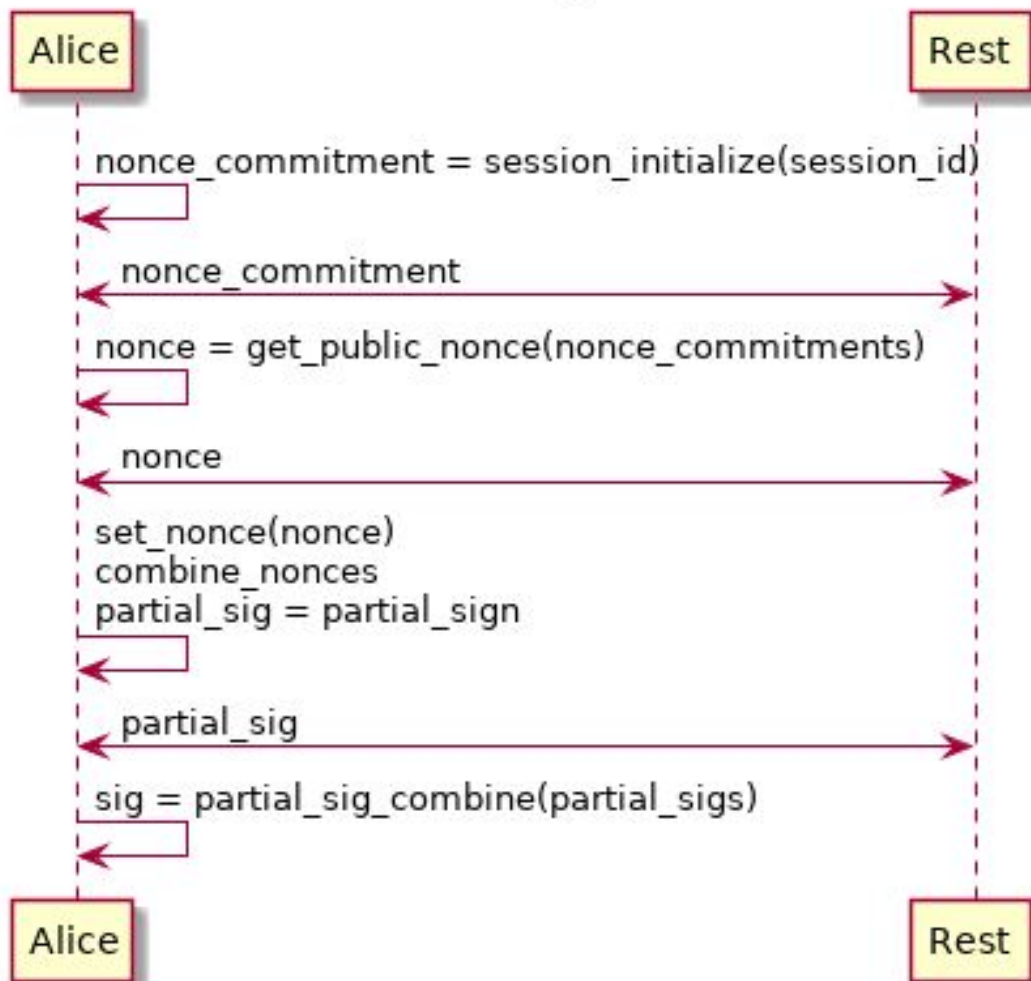
Multisignature Options with BIP-taproot

1. **use** `CHECKMULTISIG` replacement opcode `CHECKSIGADD`
 - uses BIP-schnorr and is batch verifiable
2. **Key aggregation**
 - Encode n-of-n signing policy in single public key and single BIP-schnorr signature
 - more fungible, cheaper
 - interactive protocol

Key Aggregation Options

1. **“Legacy”**: p2wpkh key aggregation
 - complicated and [80 bits security](#)
2. **BIP-taproot**: MuSig key aggregation
 - $P = \text{hash}(P1, P2, 1)P1 + \text{hash}(P1, P2, 2)P2$
3. **BIP-taproot**: Non-MuSig key aggregation
 - $P = P1 + P2$, and proof of knowledge to avoid key cancellation
 - But one party can add taproot tweak!
 - $P1 = P1' + \text{hash}(\text{prefix}, P, \text{root})G$

MuSig



MuSig Implementation

using libsecp-zkp is safe if you

1. Never reuse a session id
 - need randomness or atomic counter
2. Never copy the state
 - otherwise: Nonce reuse and active attacks

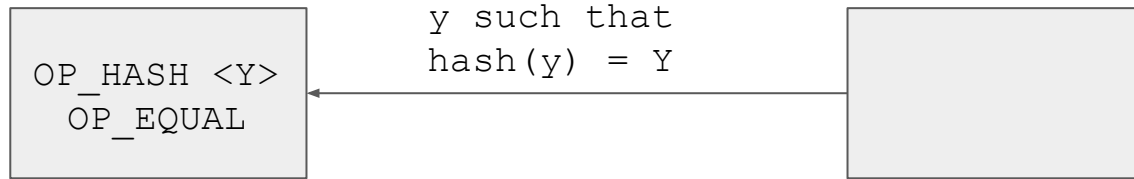
MuSig: Reducing Communication

- Can attach the nonce (commitment) to already existing messages in protocol
 - old message: `ClientHello`
 - new message: `(ClientHello, nonce_commitment)`
- Can run multiple sessions in parallel (“pre-sharing nonces”)
- Three parallel sessions get one sig per round
 - `(partial_sig_i, nonce_i+1, nonce_commitment_i+2)`

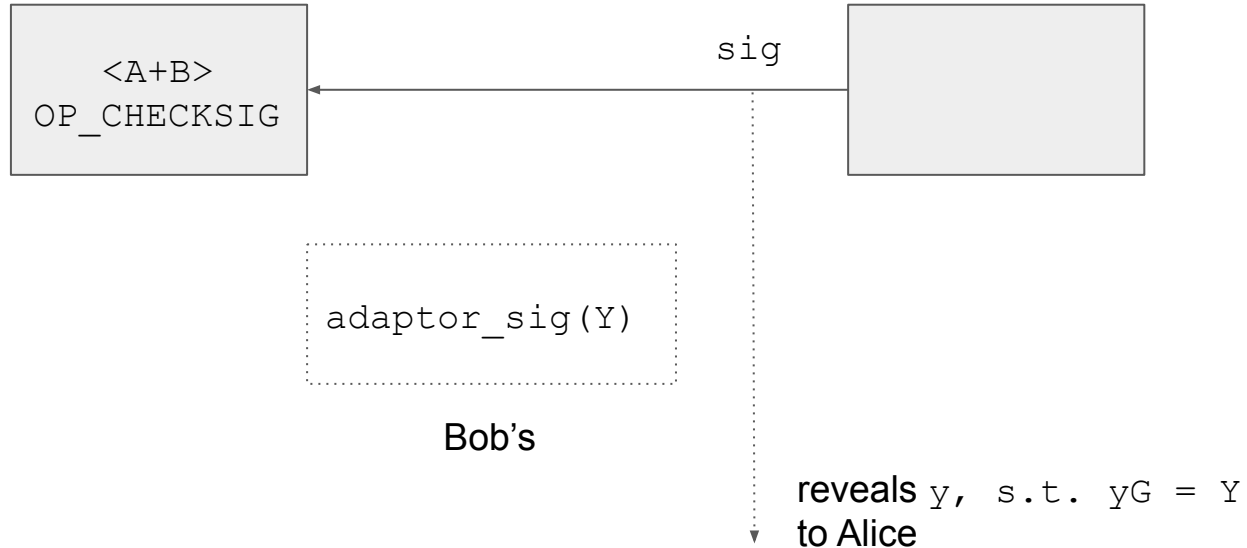
MuSig with Offline/Hardware Wallets is hard

- Storing state on persistent medium is a copy (**dangerous**)
- Therefore, serializing state not supported right now in our implementation
- Just have a “single” session?
 - Need to travel to your HWW vault for every single signature
- Hope: deterministic nonce derivation
 - no randomness, no state, two rounds
 - but must be efficient
 - adds code complexity

Hash locks



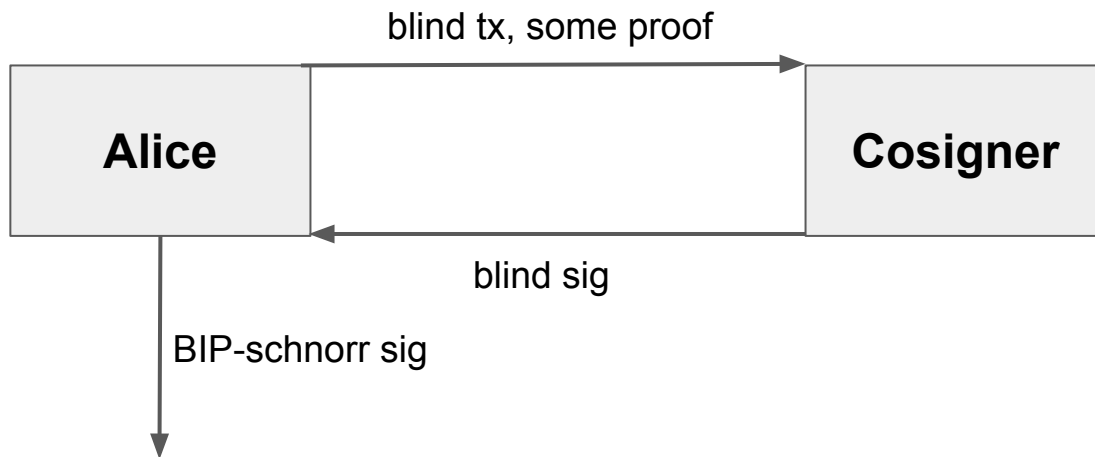
MuSig Adaptor Signatures



- **DANGER:** partial verification required
- Bonus: works with n-of-n, where $n \geq 2$

Blind Schnorr Signatures

- Interactive protocol between client and signing server
- Signer does not know the message being signed
- Result is a BIP-schnorr signature



Blind Schnorr Signatures Problems

- Vulnerable to [Wagner's attack](#)
 - 65536 parallel signing sessions can forge a signature with only $\mathcal{O}(2^{32})$ work
- Moreover, they [can't be proven secure](#) in the Random Oracle Model

Blind Schnorr Signatures

1. If you just need blind signatures (f.e. ecash)
 - Don't use blind Schnorr signatures
2. If you need blind signatures for Bitcoin transactions
 - Need to use blind Schnorr signatures
 - Idea to prevent Wagner's attack
 - i. Client blinds message with 128 different blinding factors and sends them to server
 - ii. Server picks only one of those to blindly sign

Conclusion

- BIP-taproot is a substantial efficiency & fungibility improvement
- Simple sending remains simple
- Can use libsecp256k1 ecosystem for cryptography
- DL assumption is nice (fast, studied)
 - but requires interactive protocols, creates new challenges
- TODO: k-of-n threshold signatures
- Please try to break it!
- Slides at nickler.ninja/slides/2019-breaking.pdf